

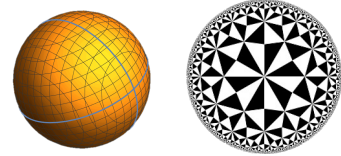
Curvature and Complexity: Better lower bounds for geodesically convex optimization

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1 Setting: g-convex optimization

$$\min_{x \in \mathcal{M}} f(x) \quad (\mathbf{P})$$

\mathcal{M} is a Riemannian manifold



f is geodesically convex (**g-convex**)

Def: For every geodesic γ , the 1D function $t \mapsto f(\gamma(t))$ is convex.

Question: complexity of **(P)** & dependence on curvature of \mathcal{M} ?

2 Curvature

$$K = 0 \qquad K < 0$$

Constant curvature:

Euclidean space \mathbb{R}^d

Hyperbolic space \mathbb{H}^d
($K = -1$)

Volume of ball of radius r :

$$\text{Vol} \sim r^d = e^{d \log(r)}$$

$$\text{Vol} \sim e^{dr}$$

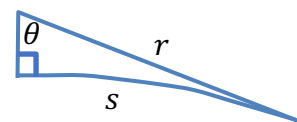
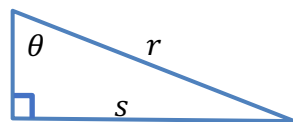
Pythagorean thm:

$$s^2 = (1 - \epsilon^2)r^2$$

$$s^2 \approx \left(1 - \frac{\epsilon^2}{\zeta}\right)r^2$$

$$\zeta = \frac{r\sqrt{|K|}}{\tanh(r\sqrt{|K|})} \sim 1 + r\sqrt{|K|}$$

$$\epsilon = \cos(\theta)$$



Applications usually have $K \leq 0$: Operator scaling, Robust covariance estimation, Matrix normal models, Intrinsic medians (e.g., for phylogenetics)

3 Complexity: the computational problem

$\mathcal{M} = \mathbb{H}^d$ is a hyperbolic space of curvature $K = -1$ (simplifying assumption)

f attains a global minimizer x^* in a known radius- r geodesic ball $B(x_{\text{ref}}, r)$

f is globally **g-convex** and regular (consider four function classes):

- (P1) Lipschitz & Low-dimensional (d fixed)
 - (P2) Lipschitz $|f(x) - f(y)| \leq M \text{dist}(x, y)$
 - (P3) Smooth $|\nabla f(x) - P_{y \rightarrow x} \nabla f(y)| \leq L \text{dist}(x, y)$
 - (P4) Smooth & strongly g-convex
- $$\left. \begin{array}{l} \text{(P1) Lipschitz \& Low-dimensional (d fixed)} \\ \text{(P2) Lipschitz } |f(x) - f(y)| \leq M \text{dist}(x, y) \end{array} \right\} f(x) - f(x^*) \leq \epsilon \cdot Mr$$
- $$\left. \begin{array}{l} \text{(P3) Smooth } |\nabla f(x) - P_{y \rightarrow x} \nabla f(y)| \leq L \text{dist}(x, y) \\ \text{(P4) Smooth \& strongly g-convex} \end{array} \right\} f(x) - f(x^*) \leq \epsilon \cdot \frac{1}{2} Lr^2$$

Black-box model: Algorithm can query an oracle at x to get $f(x), \nabla f(x)$

Worst-case complexity: Least number of oracle queries to find an ϵ -optimal point x , for all functions in the class.

- For (P2) & (P3): Depends on both ϵ and $\zeta \sim 1 + r\sqrt{|K|}$

4 Main results

g-convex setting	Lower bound	Upper bound	Algorithm
(P1) Lipschitz, lo-dim	$\tilde{\Omega}(\zeta + d)$	$O(\zeta d^2)$	Center of gravity (Rusciano'19)
(P2) Lipschitz	$\tilde{\Omega}\left(\zeta + \frac{1}{\zeta^2 \epsilon^2}\right)$	$O\left(\frac{\zeta}{\epsilon^2}\right)$	Subgradient descent (Zhang & Sra'16)
(P3) Smooth	$\tilde{\Omega}\left(\zeta + \frac{1}{\zeta \sqrt{\epsilon}}\right)$	$\tilde{O}\left(\sqrt{\zeta/\epsilon}\right)$	RNAG-C (Kim & Yang'22 / Martinez-Rubio et al.'22)
(P4) Smooth, strongly g-convex	$\tilde{\Omega}(\zeta^* + \sqrt{\kappa})$	$\tilde{O}(\sqrt{\zeta \kappa})$	RNAG-SC (Kim & Yang'22 / Martinez-Rubio et al.'22)

*Hamilton & Moitra'21 / C & Boumal'22

Note: $\kappa \geq \zeta$ and $\frac{1}{\epsilon} \geq \zeta$

Consequences:

- Provides first lower bounds depending on ϵ, κ, d
- Curvature terms in all upper bounds are unavoidable.
- "Full" acceleration is impossible for smooth g-convex opt (P3)

But, lower bounds do not match upper bounds! To address this:

- **Tight $\Omega\left(\frac{\zeta}{\epsilon^2}\right)$ lower bound for subgradient descent** (with Polyak step sizes) for (P2)
- **$\tilde{\Omega}(\zeta d)$ lower bound for cutting-planes game**, a proxy for (P1)

5 Lower bound $\tilde{\Omega}(\zeta)$ for (P1), (P2), (P3)

Slight modification of hard function in C & Boumal'22



Key geometric fact: Volume of ball of radius r scales like e^{dr}

- Pack $B(x_{\text{ref}}, r)$ with exponentially many balls $B\left(z_j, \frac{r}{5}\right)$
- Initially, each z_j is a potential value for x^*

6 Lower bound $\Omega\left(\frac{1}{\zeta^2 \epsilon^2}\right)$ for (P2)

In Euclidean space: $f(x) = \max_{i=1, \dots, d} \{ \langle s_i e_i, x \rangle \}, s_i \in \{-1, +1\}$

Fundamental difficulty: No good notion of linear functions on manifolds.

Totally geodesic submanifolds: $S_i = \{x \in \mathcal{M} = \mathbb{H}^d : \langle g_i, \log_{y_i}(x) \rangle = 0\}$

- Generalization of affine subspace

Idea: use **max of distance to totally geodesic submanifolds**:

$$f(x) = \max_{i=1, \dots, d} \{ \text{dist}(x, S_i^{s_i}) \}, s_i \in \{-1, +1\}$$

7 Lower bound $\tilde{\Omega}\left(\frac{1}{\zeta \sqrt{\epsilon}}\right)$ for (P3)

Worst function in world: $f(x) = x_{(1)}^2 - 2x_{(1)} + \sum_i (x_{(i)} - x_{(i+1)})^2 + x_{(k)}^2$

Smooth the Lipschitz lower bound with Moreau envelope (Guzman & Nemirovski'15):

$$f_\lambda(x) = \inf_{y \in \mathcal{M}} \left\{ f(y) + \frac{1}{2\lambda} \text{dist}^2(x, y) \right\}$$

Surprising because no good notion of Fenchel dual on manifolds.

8 Lower bound $\Omega\left(\frac{\zeta}{\epsilon^2}\right)$ for subgrad descent (P2)

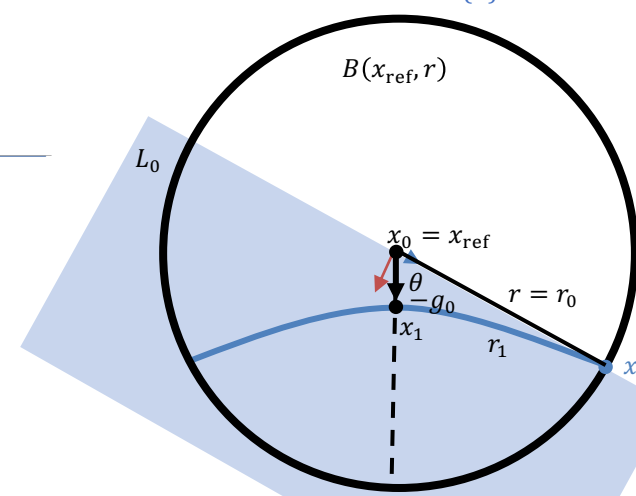
New worst function in the world (not an extension of a Euc. proof):

$$f(x) = \text{dist}(x, x^*) + \max_{i=0, \dots, d-2} \left\{ \frac{1}{4\epsilon} \text{dist}(x, L_i) \right\}$$

L_i are cleverly chosen hyperbolic halfspaces.

Key Geometric fact: Pythagorean theorem $r_1^2 \approx \left(1 - \frac{\epsilon^2}{\zeta}\right)r^2$

Construction for (8)



Construction for (6)

